# EECE.3170: Microprocessor Systems Design I 

Summer 2017

## Homework 5 Solution

For each of the following complex operations, write a sequence of PIC 16F1829 instructions that performs an equivalent operation. Assume that $X, Y$, and $Z$ are 16-bit values split into individual bytes as shown in the following cblock directive, which defines two additional variables you can use:

```
cblock 0x70
    XH, XL ; High and low bytes of X
    YH, YL ; High and low bytes of Y
    ZH, ZL ; High and low bytes of Z
    TEMP ; Temporary byte, if needed
endc
```

a. Perform the 16-bit addition: $X=Y+Z$. Do not change $Y$ or $Z$ when performing this operation.

Solution: First, we'll look at the inefficient method, which would work on any PIC microcontroller:

| movf | YL, W | ; Copy YL to XL |
| :--- | :--- | :--- |
| movwf | XL |  |
| movf | YH, W | ; Copy YH to XH |
| movwf | XH |  |
| movf | ZL, W | ; Add low bytes |
| addwf | XL, F |  |
| btfsc | STATUS, C | ; Account for carry |
| incf | XH, F |  |
| movf | ZH, W | ; Add high bytes |
| addwf | XH, F |  |

We can do this operation more efficiently by using the addwfc instruction found on microcontrollers like the PIC16F1829, which allows you to get rid of the extra instructions that "account for the carry:"

| movf | YL, W | ; Copy YL to XL |
| :--- | :--- | :--- |
| movwf | XL |  |
| movf | YH, W | ; Copy YH to XH |
| movwf | XH |  |
| movf | ZL, W | ; Add low bytes |
| addwf | XL, F |  |
| movf | ZH, W | ; Add high bytes, including carry from low byte |
| addwfc | XH, F |  |

b. Perform the 16-bit subtraction: $X=Y-Z$. Do not change $Y$ or $Z$ when performing this operation.

Solution: This operation is very similar to 16-bit addition, although you have to be more careful about what register is moved into the working register before the subtract instructions. First, the inefficient version-remember that $\mathrm{C}=0$ if a borrow occurs:

| movf | YL, W | ; Copy YL to XL |
| :--- | :--- | :--- |
| movwf | XL |  |
| movf | $\mathrm{YH}, \mathrm{W}$ | ; Copy YH to XH |
| movwf | XH |  |
| movf | $\mathrm{ZL}, \mathrm{W}$ | ; Subtract low bytes |
| subwf | $\mathrm{XL}, \mathrm{F}$ |  |
| btfss | STATUS, C | ; Account for borrow $(\mathrm{C}=0 \rightarrow$ "borrow" = 1) |
| decf | XH, F |  |
| movf | $\mathrm{ZH}, \mathrm{W}$ | ; Subtract high bytes |
| subwf | XH, F |  |

And the more efficient version that uses the "subtract with borrow" subwfb instruction:

| movf | $\mathrm{YL}, \mathrm{W}$ | ; Copy YL to XL |
| :--- | :--- | :--- |
| movwf | XL |  |
| movf | $\mathrm{YH}, \mathrm{W}$ | ; Copy YH to XH |
| movwf | XH |  |
| movf | $\mathrm{ZL}, \mathrm{W}$ | ; Subtract low bytes |
| subwf | $\mathrm{XL} F$, |  |
| movf | $\mathrm{ZH}, \mathrm{W}$ | ; Subtract high bytes, taking borrow into |
| subwfb | $\mathrm{XH}, \mathrm{F}$ | ; account |

c. Perform a 16-bit arithmetic right shift: $X=Y \gg Z L$. (Note that, because the shift amount is no greater than 15, a single byte is sufficient to hold that value.) Do not change $Y$ or $Z L$ when performing this operation.

Solution: Similarly to the last two problems, the first thing to be done is move the value to be shifted into the destination registers XH and XL. Once that's done, set up a loop with ZL iterations (we'll have to copy that value to another register so it's not changed) and do the shift. Remember, while the shift for the high byte can be an arithmetic shift, we need a rotate instruction to change the low byte so that the bit shifting between bytes is correctly accounted for.

|  | movf | YL, W | ; Copy YL to XL |
| :---: | :---: | :---: | :---: |
|  | movwf | XL |  |
|  | movf | YH, W | ; Copy YH to XH |
|  | movwf | XH |  |
|  | movf | ZL, W | ; Copy ZL to TEMP |
|  | movwf | TEMP |  |
| L: | asrf | XH, F | ; Shift upper byte ( $\mathrm{C}=$ bit to be shifted into XL ) |
|  | rrf | XL, F | ; Shift lower byte |
|  | decfsz | TEMP, F | ; Decrement loop counter and return to start |
|  | goto | L | ; of loop if there are more iterations. |

d. Given an 8-bit variable, YL, perform the multiplication:

$$
Y L=Y L * 10
$$

Hint: Note that multiplication by a constant amount can be broken into a series of shift and add operations. For example, in general:

- $\quad X^{*} 2$ can be implemented by shifting $X$ to the left by $1(X \ll 1)$
- $X * 5$ can be implemented as $(X * 4)+X=(X \ll 2)+X$

Solution: Recognize that $\mathrm{YL} * 10=(\mathrm{YL} * 8)+(\mathrm{YL} * 2)=(\mathrm{YL} \ll 3)+\mathrm{YL}+\mathrm{YL}$

|  | movf | YL, W | ; Copy original value of YL into TEMP |
| :---: | :---: | :---: | :---: |
|  | movwf | TEMP |  |
|  | movlw | 3 | Set W = 3—use as loop counter for left shift |
| L: | Islf | YL, F |  |
|  | addlw | -1 | ; Decrement loop counter |
|  | btfss | STATUS, Z | ; and exit once it reaches 0 |
|  | goto | L |  |
|  | movf | TEMP, W | ; $\mathrm{W}=\mathrm{TEMP}=$ original value of YL |
|  |  |  | $\begin{aligned} & \text {; At this point, } \mathrm{YL}=(\text { original } \mathrm{YL}) \ll 3 \\ & ;=(\text { original } \mathrm{YL}) * 8 \end{aligned}$ |
|  | addwf | YL, F | ; $\mathrm{YL}=$ (original YL ) * 9 |
|  | addwf | YL, F | ; $\mathrm{YL}=($ original YL$) * 10$ |

e. Given two 8 -bit variables stored in XL and YL, copy the value of bit position YL within variable XL into the carry flag. For example:

- If $X L=0 x 03$ and $Y L=0 x 00$, set $C$ to the value of bit 0 within $X L$.
- Since $X L=0 x 03=0000001 \underline{1}_{2}, C=1$
- If $X L=0 x C 2$ and $Y L=0 x 04$, set $C$ to the value of bit 4 within $X L$.
- Since $X L=0 x C 2=11000011_{2}, C=0$

Note that:

- This operation is very similar to the bit test (BT) instruction in the $x 86$ architecture.
- Since YL is not a constant, you cannot use the value of YL directly in any of the PIC bit test instructions (for example, btfsc XL, YL is not a valid instruction).
- Your code should not modify either XL or YL.


## Solution

| movlw | $0 \times 01$ | ; TEMP will hold bit mask used |
| :---: | :---: | :---: |
| movwf | TEMP | to isolate bit YL within XL |
| movf | YL, W | ; Copy B to W-determines \# of times to shift temp |
| btfsc | STATUS, Z | ; Once W hits 0, end loop-bit mask is set |
| goto | L2 | ; Must test this first for case where YL == 0 |
| rlf | TEMP, F | ; TEMP will eventually be $1 \ll \mathrm{YL}$ |
| addlw | -1 | ; Decrement W |
| goto | L |  |
| bcf | STATUS, C | ; Clear C |
| movf | TEMP, W | ; AND temp with XL to mask out all but bit YL |
| andwf | XL, W |  |
| btfss | STATUS, Z | ; If result is non-zero, set C bit; otherwise, |
| bsf | STATUS, C | ; leave as 0 |

