# 16.482 / 16.561: Computer Architecture and Design <br> Spring 2015 

Homework \#9
Due Thursday, 4/16/15

## Notes:

- While typed submissions are preferred, handwritten submissions are acceptable.
- Any electronic submission must be in a single file. Archive files will not be accepted.
- Electronic submissions should be e-mailed to Dr. Geiger at Michael_Geiger@uml.edu.
- This assignment is worth a total of 100 points.

1. RAID ( 40 points) You are working with a 5-disk RAID array that contains a total of 15 sectors; the exact sector configuration depends on the RAID level used. In all cases, twelve of the fifteen sectors (S0-S11) will hold data, while the remaining three sectors (P0-P2) hold parity information. Large reads and writes (reads/writes that access an entire stripe in the array) take 300 ms , small reads (reads involving only a single disk) take 150 ms , and small writes (writes involving 1 data disk +1 parity disk) take 200 ms .

Given the following sequence of sector reads and writes, determine the time required if the array is configured with RAID 3, RAID 4, and RAID 5. Assume the following:

- Requests are queued in such a manner that two consecutive operations may proceed simultaneously if they do not share any disk within the array.
- If two disks, $D_{x}$ and $D_{y}$, are in use, and the access to $D_{x}$ finishes before the access to $D_{y}$, a new operation may start immediately assuming it does not involve $D_{y}$.
- Multiple accesses to the same stripe may overlap if they don't use the same disk.

1. read S 0
2. write S 5
3. write S 8
4. read S9
5. read S3
6. write S7
7. read S11
8. write S1

In each case, show the organization of the array to support your answer.
2. Snooping protocols ( 30 points) You are given a four-processor system that uses a writeinvalidate, snooping coherence protocol. Each direct-mapped, write-back cache has four lines, each of which holds eight bytes; in the diagram below, only the least-significant byte of each word is shown. The cache states are I (invalid), S (shared), and M (modified/exclusive). The caches and memory have the following initial state:

|  | State | Tag | Data |  |
| :---: | :---: | :---: | :---: | :---: |
| B0 | 1 | 0x100 | 01 | 23 |
| B1 | S | 0x108 | 00 | 88 |
| B2 | M | 0x110 | 00 | 30 |
| B3 | 1 | 0x118 | 00 | 10 |


|  | State | Tag | Data |  |
| :---: | :---: | :---: | :---: | :---: |
| B0 | I | 0x100 | 01 | 23 |
| B1 | M | 0x128 | AB | CD |
| B2 | S | 0x130 | 14 | 12 |
| B3 | S | 0x118 | 14 | 92 |


| P2 | State | Tag | Data |  |
| :---: | :---: | :---: | :---: | :---: |
| B0 | S | 0x120 | 13 | 31 |
| B1 | S | 0x108 | 00 | 88 |
| B2 | I | 0x130 | 51 | 55 |
| B3 | I | 0x138 | 01 | 38 |

## P3

|  | State |  | Tag |  |
| :---: | :---: | :---: | :---: | :---: |
|  | Data |  |  |  |
| B0 | $0 \times 120$ | 13 | 31 |  |
|  | S | $0 \times 108$ | 00 | 88 |
|  | B2 | I | $0 \times 110$ | 00 |


| Memory |  |  |
| :---: | :---: | :---: |
| Address | Dat |  |
| 0x100 | 00 | 00 |
| 0x108 | 00 | 88 |
| $0 \times 110$ | 20 | 08 |
| $0 \times 118$ | 14 | 92 |
| 0x120 | 13 | 31 |
| 0x128 | FF | FE |
| 0x130 | 14 | 12 |
| 0x138 | AB | BA |

For each of the transactions listed below, list all cache blocks modified and their final state, as well as all memory blocks modified and their final state. Assume each set of transactions starts with the same initial state-in other words, your answer to part (b) does not depend on your answer to part (a). However, you should track the state transitions of each block throughout the problem.
a. $\quad \mathrm{P} 1:$ read $0 \times 128$
P2: read 0x128
c. $\quad$ P3: write $0 \times 128 \leftarrow 45$
P0: write $0 \times 128 \leftarrow 67$
P3: read $0 \times 12 \mathrm{C}$
P3: read 0x128
b. P2: write $0 \times 110 \leftarrow 14$

P2: write 0x114 < DF
P2: write 0 x130 $\leftarrow 20$
3. Directory protocols ( 30 points) Say we have a four-processor system that uses a writeinvalidate, directory coherence protocol. The system contains a total of 8 memory blocks, as shown in the initial directory state below:

| Block \# | P0 | P1 | P2 | P3 | Dirty |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 |
| 2 | 1 | 1 | 1 | 1 | 0 |
| 3 | 0 | 1 | 1 | 0 | 0 |
| 4 | 0 | 0 | 0 | 1 | 1 |
| 5 | 0 | 0 | 0 | 0 | 0 |
| 6 | 0 | 1 | 0 | 0 | 0 |
| 7 | 1 | 0 | 0 | 0 | 1 |
| 8 | 0 | 0 | 0 | 0 | 0 |

a. (5 points) What state is each block in?
b. (25 points) What messages are sent between the nodes and directory for each of the following transactions to ensure the processor has the most up-to-date block copy, the appropriate invalidations are made, and the directory holds the appropriate state?
i. P1: write block 0
ii. P2: read block 7
iii. P3: write block 4
iv. P0: write block 2
v. P1: read block 6

